

Problem 2. (15 points)

Let \mathcal{N} be the set of natural numbers (i.e. positive integers). Let $f : \mathcal{N} \rightarrow \mathcal{N} \cup \{0\}$ be defined as follows. For $n \in \mathcal{N}$, let $f(n)$ be the largest non-negative integer p for which 2^p divides n . So $f(1) = 0, f(2) = 1, f(3) = 0, f(4) = 2, f(5) = 0, f(6) = 1, f(7) = 0, f(8) = 3$, etcetera. Now define a binary relation g on \mathcal{N} as follows:

$$g : \mathcal{N} \times \mathcal{N} \rightarrow \{\text{true}, \text{false}\},$$

$$g(m, n) = \begin{cases} \text{true} & \text{if } f(m) - f(n) \text{ is divisible by 3} \\ \text{false} & \text{otherwise} \end{cases} .$$

Carefully show that g is an equivalence relation on \mathcal{N} . *Hint:* This looks more complicated than it really is.

Problem 3. (15 points)

We know that the union of regular language is regular. Using this fact, give a *careful proof by induction* that every finite language is regular. To help make sure you do this properly, I recommend that you use the notation $P(n)$ to mean that “any language of finite size n is regular”. Don’t forget about the empty language.

Problem 4. (16 points)

Determine the language of each of following **NFAs**. Don't be too brief or too verbose. Make sure your meaning is clear. Think about your wording (and perhaps notation) before you write.

(a)

(b)

(c)

(d)

Problem 5. (16 points)

Draw a **DFA** for each of the following languages. Remember that two transitions must leave each state.

(a) Binary strings containing 101 as a substring.

(b) $\{\epsilon, 0, 00, 1\}$

(c) ϕ

(d) Binary strings of length at least two whose next to last symbol is a zero.

Problem 6. (18 points)

Use the “subset construction” to produce a **DFA** that is equivalent to the **NFA** in part (a) of Problem 4. Do not short-cut this in any way. Do the whole construction. Remember that you should have sixteen states and two transitions out of each of these states.